A Tiled-Table Convention for Compressing FITS Binary Tables

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1 Overview

This document describes a convention for compressing FITS binary tables that is modeled after the FITS tiled-image compression method (White et al. 2009) that has been in use for about a decade. The input table is first optionally subdivided into tiles, each containing an equal number of rows, then every column of data within each tile is compressed and stored as a variable-length array of bytes in the output FITS binary table. All the header keywords from the input table are copied to the header of the output table and remain uncompressed for efficient access. The output compressed table contains the same number and order of columns as in the input uncompressed binary table. There is one row in the output table corresponding to each tile of rows in the input table. In principle, each column of data can be compressed using a different algorithm that is optimized for the type of data within that column, however in the prototype implementation described here, the gzip algorithm is used to compress every column.

When compressing most FITS tables, a lossless algorithm will be required in order to exactly preserve the values in each column. In principle, however, a lossy compression algorithm could be used to achieve higher compression in cases where the values in a particular column do not need to be exactly preserved.

This convention currently can only be used to compress FITS binary tables and is not applicable to FITS ASCII tables.

2 Algorithm Details

The algorithm for compressing a FITS binary table consists of the following sequence of steps:

A. **Divide Table into Tiles (Optional)**

In order to limit the amount of data that must be manipulated at one time, large FITS tables may be optionally divided into tiles, where each tile contains the same number of rows from the input table, except for the last tile which may contain fewer rows. Each tile is compressed in turn, and is stored in a row in the output compressed table. This convention places no restriction on the tile size, but in practice, it is recommended that FITS tables that are larger than about 10 MB in size should be divided into tiles so as to not impose too large of a memory resource burden on software that must uncompress the table.

The number of rows of data in each tile is defined by the `ZTILELEN` keyword, as described in Section 3. If this keyword is not present, then it should be assumed that the entire table has been compressed as a single tile.

B. **Transpose the Rows and Columns**

The data cells in a FITS binary table are natively stored in row by row order, where the values in the first row for each column are given in order, followed by the values for the second row, and so on. Because binary tables can contain very heterogeneous types of data in different columns, it can be difficult to directly compress this native stream of data values. One can almost always obtain better compression by internally transposing the table values into 'column-major' order, where all the values for the first column occur first, followed by all the values for the second column, and so on. Then each column of values can be compressed separately, and, at least in principle, a different compression algorithm may be chosen for each column that is optimized for that particular type of data.

Note that fixed-length vector columns are transposed in the same way as scalar columns. For example, if the table contains a `TFORMn = '30I'` column (an array of 30 16-bit integers) and contains 50 rows, then the transposed column will contain an array of $30 \times 50 = 1500$ 16-bit integers.

C. **Compress each Column**

Each column of data is compressed with a suitable compression algorithm. If the table has been divided into tiles, then the same compression algorithm must be used for a given column in all the tiles. The mnemonic name of the compression algorithm that is used for each column is given by the corresponding `ZCTYPn` keyword, as described in Section 3.

D. **Store the Compressed Bytes**

The compressed stream of bytes for each column is stored in the corresponding column in the output table. The compressed table will have exactly the same number and order of columns as in the input table, however the data type of the columns in the output table will all have a variable-length byte data type, with TFORMn = '1PB', which is appropriate for storing the compressed byte stream. Each row in the output compressed table corresponds to a tile of rows in the input uncompressed table. If the input table is compressed as a single tile, then the output table will only contain one row.

3 Keywords

With only a few exceptions, all the keywords from the uncompressed table are copied verbatim, in order, into the header of the compressed table. The header keywords remain uncompressed in order to preserve fast read and write access. Note in particular that the values of the reserved column descriptor keywords `TTYPEn`, `TUNITn`, `TSCALn`, `TZEROn`, `TNULLn`, `TDISPn`, and `TDIMn`, as well as all the column-specific WCS keywords defined in the FITS standard, have the same values in both the original and in the compressed table, with the understanding that these keywords apply to the values in the *uncompressed columns*.

The only keywords that are not copied verbatim from the uncompressed table header to the compressed table header are the mandatory `NAXIS1`, `NAXIS2`, `PCOUNT`, and `TFORMn` keywords (and possibly the reserved `THEAP` keyword), which must necessarily describe the actual structure of the compressed table. The original values of these keywords are stored in a new set of reserved keywords in the compressed table header. The complete set of keywords that have a reserved meaning within the header of a tile-compressed binary table are listed below:

- `ZTABLE` (required keyword). The value field of this keyword shall contain the logical value T. It indicates that the FITS binary table extension contains a tile-compressed binary table.
- `ZTILELEN` (optional keyword). The value of this keyword shall contain an integer representing the number of rows of data from the original binary table that are contained in each tile of

compressed data. The number of rows in the last tile, which may be smaller than the previous tiles, is given by ZNAXIS2 - ((NAXIS2 - 1) $\times$ ZTILELEN). If the ZTILELEN keyword is absent, then it should be assumed that the whole table is compressed as a single tile; the compressed table must only have a single row in this case.

- ZCTYPn (optional indexed keywords). The value field of these keywords shall contain a character string giving the mnemonic name of the algorithm that is used to compress and decompress column n of the table. The current allowed values are GZIP_1, GZIP_2, and RICE_1, and the corresponding algorithms are described in Section 4. If this keyword is absent for a given column, then the default value of GZIP_2 should be assumed.
- ZNAXIS1 (required keyword). The value field of this keyword shall contain an integer that gives the value of the NAXIS1 keyword in the original FITS table header. This represents the width in bytes of each row in the uncompressed table.
- ZNAXIS2 (required keyword). The value field of this keyword shall contain an integer that gives the value of the NAXIS2 keyword in the original FITS table header. This represents the number of rows in the uncompressed table.
- ZPCOUNT (required keyword). The value field of this keyword shall contain an integer that gives the value of the PCOUNT keyword in the original FITS table header.
- ZFORMn (required indexed keywords). These required array keywords supply the character string value of the corresponding TFORMn keyword in the original FITS table header.
- ZTHEAP (optional keyword). The value field of this keyword shall contain an integer that gives the value of the THEAP keyword, if it exists, in the original FITS table header. In practice, this keyword is rarely used.
- ZHEAPPTR (optional keyword). The value field of this keyword shall contain an integer that gives the byte offset from the start of the main data table to the beginning of the compressed heap. Usage of this keyword is described in Section 5.2.

4 Supported Compression Algorithms

This section describes the currently defined compression algorithms. Other compression algorithms may be added in the future.

4.1 GZIP_1 compression algorithm

This compression algorithm is designated by the keyword ZCTYPn = 'GZIP_1'. Gzip is the compression algorithm used in the widely distributed GNU free software utility of the same name. It was created by Jean-loup Gailly and Mark Adler. It is based on the DEFLATE algorithm, which is a combination of LZ77 and Huffman coding. Further information about this compression technique is readily available online on the Web. The "gzip -1" option is generally used which significantly improves the compression speed with only a small loss of compression efficiency.

4.2 GZIP_2 compression algorithm (default)

This compression algorithm is designated by the keyword `ZCTYPn = 'GZIP_2'`. If the `ZCTYPn` keyword is absent for a given column, then this algorithm should be assumed. This algorithm is a variation of the `GZIP_1` algorithm in which the bytes in the arrays of numeric data columns are preprocessed by shuffling them so that they are arranged in order of decreasing significance before being compressed. For example, a 5-element array of 2-byte (16-bit) integer values, with an original byte order of

```
A1 A2 B1 B2 C1 C2 D1 D2 E1 E2,
```

will have the following byte order after shuffling the bytes:

```
A1 B1 C1 D1 E1 A2 B2 C2 D2 E2.
```

Byte shuffling is only performed for numeric binary table columns that have `TFORMn` data type codes of `I, J, K, E, D, C,` or `M`. The descriptors for all variable-length array columns, with a `P` or `Q` type code, are also byte-shuffled. The bytes in columns that have a `L, X,` or `A` type code are not shuffled.

This byte-shuffling technique has been shown to be especially beneficial when compressing floating-point values because the bytes containing the exponent and the most significant bits of the mantissa are often similar for all the floating point values in the array. Thus these repetitive byte values generally compress very well when grouped together in this way. HDF Group has used this byte-shuffling technique when compressing HDF5 data files (HDF 2000).

4.3 RICE_1 compression algorithm

This compression algorithm is designated by the keyword `ZCTYPn = 'RICE_1'`. The Rice algorithm (Rice, R. F., Yeh, P.-S., and Miller, W. H. 1993, in Proc. of the 9th AIAA Computing in Aerospace Conf., AIAA-93-4541-CP) is very simple and fast. It requires only enough memory to hold a single block of 32 integers at a time and is able to adapt very quickly to changes in the input array statistics. This algorithm can only be used to compress integer table columns. Note that the byte-shuffling technique that is used with the GZIP_2 algorithm must not be applied to the integer columns when using the Rice compression algorithm..

5 Compressing Variable-Length Array Columns

Compression of binary tables that contain variable-length array columns, with a `P` or `Q` data type code, requires special consideration because the table elements are not stored in the table directly, but instead are stored in what is called the 'data heap' which follows the main table. The columns in the main data table itself contains a 'descriptor', which is composed of 2 integers that give the size and location of the array in the heap. Thus, in addition to compressing the descriptors values, one must also compress the actual array values in the heap.

5.1 Compression of the Descriptors

The descriptor values in the main table for `P` or `Q` type variable-length array columns consist of a pair of 32-bit integers or 64-bit integers, respectively. For compression purposes, these columns are treated the same as a column with a `TFORMn = '2J'` or `TFORMn = '2K'` data type.

5.2 Heap compression

After the main data table has been compressed, the entire data heap in the input table is compressed as a single block of data using the GZIP_2 algorithm (described in Section 4.2). The size of the data heap is given by the `PCOUNT` keyword (in the input table) and it begins with the byte immediately following the main data table. (In fact, the `THEAP` keyword in the input table can be used to leave a gap between the end of the main data table and the beginning of the actual heap, but this technicality can be safely ignored here). To improve the compression efficiency, the bytes in each numeric variable-length array within the heap are shuffled so that the most significant byte of each array element is given first, followed by the next most significant byte of every element, and so on. This can be done efficiently by transversing every variable-length array column in the input table that has a `I, J, K, E, D, C,` or `M` data type, in order from the first row to the last. (When uncompressing the heap, this byte unshuffling should be done in reverse order to ensure that any overlapping arrays in the heap are restored to their original byte order). The byte-shuffled heap is then compressed with the gzip algorithm, and the resulting compressed stream of bytes is stored in the heap of the output compressed table *immediately following* the previously written compressed bytes from the main data table. The starting byte location of the compressed heap, relative to the starting location of the main data table is given by the reserved `ZHEAPPTR` keyword.

The value of the `PCOUNT` keyword in the output table will give the total size of the heap in the compressed table, which includes the size of the compressed columns of data from the main data table plus the size of the compressed heap.

6 Prototype Implementation

The advantages of tiled-table compression have been demonstrated on a small sample of astronomical FITS binary tables, mostly from the HEASARC archive, using a prototype version of the CFITSIO library that supports this compression method. (Tests on a wider sample of tables are pending). Each table was compressed in three different ways. In the first test, each FITS table was simply compressed using the gzip utility program. This is currently the most commonly used method for compressing FITS tables. For comparison, the second test used the GZIP_1 compression method where the rows and columns in the table are transposed before compressing each column using gzip. Finally, the third test used the GZIP_2 method where the bytes in all the numeric columns are shuffled into decreasing order of significance, in addition to transposing the rows and columns. The following table shows the compression ratio (original table size divided by the compressed table size) for these 3 different compression methods. The 13 different FITS tables have been sorted in order of increasing gzip compression ratio.

File	gzip	GZIP_1	GZIP_2	Disk Savings factor
1	1.11	1.16	1.49	3.3
2	1.41	1.72	2.02	1.7
3	1.48	1.70	1.94	1.5
4	1.51	1.81	2.12	1.6
5	1.61	1.95	2.34	1.5
6	1.71	1.89	1.95	1.2
7	1.73	2.43	2.57	1.4
8	1.83	2.14	3.08	1.5
9	2.63	2.63	3.75	1.2
10	2.88	3.93	6.35	1.3
11	3.07	4.88	3.13	1.0
12	3.85	5.51	22.59	1.3
13	5.39	7.63	7.60	1.1

The amount of compression that is achieved depends greatly on the particulars of the data contained in the table. However, in nearly every case, transposing the rows and columns in the table (using GZIP_1) significantly improves the compression ratio, and shuffling the bytes in the numeric columns (using GZIP_2) provides even further improvement. The only real exception is file 11, where shuffling the bytes results in worse compression. This particular table is unusual, however, because the data values in most of the columns are nearly the same in every row. Gzip is apparently better able to compress these repetitive values if the individual bytes are not shuffled.

The last column in the table gives the ratio of the amount of disk space that is saved when compressing the table using the GZIP_2 method versus just using the gzip utility program to compress the whole table. On average, GZIP_2 saves 1.5 times more disk space.

Using this current prototype implementation, it takes about 50% more CPU time to compress the table using the GZIP_2 method than when using the gzip utility. It is expected, however, that further optimization of the code will eliminate most of this time difference.

Finally, it should be noted that these tests were performed on large FITS tables where the size of the required FITS primary array header, plus the size of the table header, plus the size of any fill bytes at the end of the last FITS 2880-byte block of the table, is insignificant compared to the size of the data table itself. Because this tiled-table compression method does not compress the headers, and because the minimum size of the compressed table is 1 FITS block (2880 bytes), this compression method will be less effective for small tables or for table where the headers account for a significant fraction of the size of the table.

We anticipate that further prototyping work on this table compression convention will be performed in the near future. In particular, the performance of other lossless compression algorithms beside gzip should to be explored. It may also be desirable to consider extending this convention to support compression of ASCII FITS tables, and possibly add an option for compressing the headers of images or tables that contain a relatively large number of keywords.

References

HDF 2000, "Performance Evaluation Report: gzip, bzip2 compression with and without shuffling," `http://www.hdfgroup.org/HDF5/doc_resource/H5Shuffle_Perf.pdf` (not to be confused with `http://www.youtube.com/watch?v=YjfHRRrFUPk`)

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